

# Conclusions

In this thesis we developed algorithms for the high-dimensional  $L_\infty$ -nearest-neighbor problem. We presented their average-case analysis under the assumption that the  $n$  data points are chosen independently at random under uniform distribution. We did not assume any distribution of the query point. The algorithms improve considerably the brute-force method, are simple and easy to implement, and have a low storage requirement.

The query algorithms developed in Section 3.1 are based on simple preprocessed data structures. The searching algorithm REJECT\_SCAN requires  $O(nd)$  storage and  $O(nd \ln n)$  preprocessing time, and has an expected running time of  $O\left(n \ln\left(\frac{d}{\ln n} + 1\right) + n + d \ln d\right)$ . It improves the brute-force method by a factor of  $\Omega\left(\frac{\ln(d/\ln n + 1)}{d}\right)$ . The constant in the  $O$ -term is close to one, and for many applications the dimension  $d$  is in the range of some hundreds, so the speedup factor comes close to 50, which is a considerable improvement in practice. This consideration is confirmed by the experimental comparison with the brute-force method in Section 4.5.

The searching algorithm REJECT uses the same data structure like the algorithm REJECT\_SCAN, has an expected running time of  $O(n \ln d + n + d \ln d)$  and can be extended to work efficiently in the external-memory model of computation.

The data structure based on a preprocessed partition of the data set into monotone sequences achieves an expected runtime that is sublinear in  $n$  and  $d$ . The query algorithm has an expected running time of  $O\left(\sqrt{d} \cdot n^{1-\frac{1}{\sqrt{d}}} \cdot \ln n\right)$  for dimensions  $d < \left(\frac{\ln n}{\ln \ln n}\right)^2$ .

We extended our methods to solve the  $k$ -nearest-neighbor problem with an expected asymptotic runtime of  $O\left(n \ln\left(\frac{d}{\ln(n/k)} + 1\right) + n + dk + d \log d\right)$ . Furthermore, the expected runtime analysis of our algorithms has been generalized to other "well-behaved" probability distributions.

In Chapter 5 we developed a method which provides tradeoffs between the space complexity of the data structure and the time complexity of the query algorithm.

